

Approssimazioni

$\langle M, S \rangle$

\uparrow
 L_H

$H \rightarrow S$

sp. top. synth

due tipi di famile chiuse

$\exists \in C$

τ prima di sotto

$H^m \times S^m \rightarrow S$

k -tuple
di termini

$C \subseteq S^k$

$\tau(x; u)$

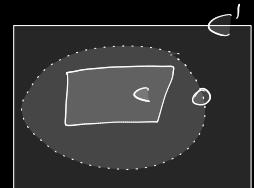
$\varphi' > \varphi$ se le occorrenze di $\exists \in C$

in φ sono sostituite da $\exists \in C'$

dove $C' > C$

OSS 1

$\varphi \rightarrow \varphi'$



OSS 2

$\exists \varphi \in L_H \quad \varphi \geq \varphi$

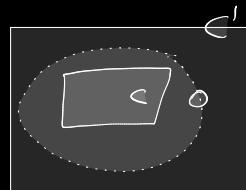
OSS 3

$\varphi' > \varphi$ esiste φ'' t.c. $\varphi' > \varphi'' > \varphi$

negazione forte

$\widehat{\varphi} \perp \varphi$ se le occorrenze di $\exists \in C$

sono sostituite da $\exists \in \widehat{C}$ dove $\widehat{C} \cap C = \emptyset$



$\exists \varphi \rightarrow \exists \varphi'$

$\forall \rightarrow \exists$

$\wedge \rightarrow \vee$

$\exists \rightarrow \forall$

$\vee \rightarrow \wedge$

$\varphi \rightarrow \neg \widehat{\varphi}$

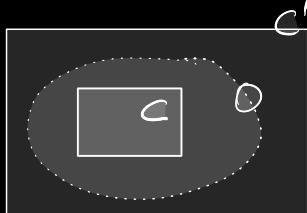
$\widehat{\varphi} \rightarrow \neg \varphi$

Prosto ① $\varphi' > \varphi$ esiste $\widehat{\varphi} \perp \varphi$ t.c. $\varphi \rightarrow \neg \widehat{\varphi} \rightarrow \varphi'$

② $\widehat{\varphi} \perp \varphi$ esiste $\varphi' > \varphi$ t.c. $\varphi \rightarrow \varphi' \rightarrow \neg \widehat{\varphi}$

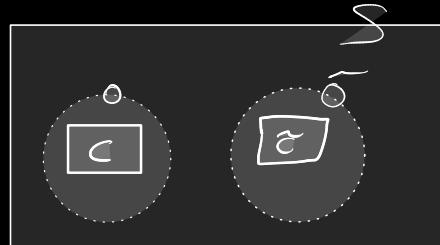
$C' = S \cdot \widetilde{S}$

Dim ①



$\widetilde{C} = C' \cdot O$

Dim ②



$\mathcal{F}^C \subseteq \mathcal{F}^P$

~~\exists_H~~

Lemma

N p-w-sistema $\varphi \in \mathcal{F}^P$

$N \models \{\varphi\}' \hookrightarrow \varphi$

dove $\{\varphi'\} = \{\varphi' : \varphi' > \varphi\}$

Dim $\varphi = (\exists \in C)$

$\{\varphi\}' = \{\exists \in C' : C' > C\} \leftrightarrow \exists \in \bigcap_{C' > C} C' = C$

possò induuttivo per esistenziale:

$\{\varphi(x)\}' \xrightarrow{I_H} \varphi(x)$

$\{\exists \varphi(x)\}' \xrightarrow{OBIETTIVO} \exists \varphi(x)$

Riposo $T \perp M$

in un modello sufficientemente saturo

$$\exists y P(x, y) = \{\exists y \varphi(x, y) : \varphi(x, y) \in P\}$$

Controesempio \models $S = [0, 1]$ $M = [0, 1]$ $i: [0, 1] \rightarrow [0, 1]$

$L_4 = \{x\}$ M ideale \subseteq

$$\varphi = \exists x \left[\underbrace{x \neq 0} \wedge i(x) \in \{0\} \right]$$

per ogni $\varphi' > \varphi$

$$M \models \neg \varphi \quad M \models \varphi'$$

M è ρ -modello se $\exists N \models^{\rho} M \ V \models \varphi \Rightarrow M \models \varphi$
 se M ρ -modello $M \models^{\rho} N \Leftrightarrow (M \models \varphi \Leftrightarrow N \models \varphi)$

Lemma LSASE $\begin{array}{l} \textcircled{1} M \text{ } \rho\text{-modello} \\ \textcircled{2} M \models \{\varphi\}' \rightarrow \varphi \end{array}$

Dim $\textcircled{1} \Rightarrow \textcircled{2}$ $M \models \{\varphi\}'$ prendere $N \models M$ ρ -satura
 $N \models \{\varphi\}'$ quindi $N \models \varphi$ quindi $M \models \varphi$.
 $\textcircled{2} \Rightarrow \textcircled{1}$ $M \models^{\rho} N \models \varphi$ basta mostrare che $M \models \{\varphi\}'$ per questo
 prendono $\varphi' > \varphi$ t.c. $M \models \neg \varphi'$ esiste $\widehat{\varphi} \perp \varphi$ t.c. $\varphi \rightarrow \neg \widehat{\varphi} \rightarrow \varphi'$
 $M \models \widehat{\varphi}$ quindi $N \models \widehat{\varphi}$

\mathcal{U} ρ -satura modello $M \models^{\rho} \mathcal{U}$
 $M \models \varphi \Rightarrow \mathcal{U} \models \varphi$

Fatto M modello
 $M \models \{\varphi\}' \Leftrightarrow \mathcal{U} \models \{\varphi\}'$

esiste $\varphi(x) \in P$

Fatto $P(x) \subseteq \mathcal{F}$ $\textcircled{1}$ se $P(x) \rightarrow \neg \varphi(x)$ allora $\sqrt{\neg \varphi(x)} \rightarrow \neg \varphi(x)$
 $\varphi(x) \in \mathcal{F}$ $\textcircled{2}$ se $P(x) \rightarrow \varphi(x)$ per ogni $\varphi' > \varphi$ esiste $\varphi(x) \in P$
 $\varphi(x) \rightarrow \varphi'(x)$

Dimo $\textcircled{1}$ $P(x) \cup \{\varphi(x)\} \vdash \perp$ ---- contradd.

Dim $\varphi(\alpha) \cup \{\neg\varphi(\alpha)\} \vdash \perp$ STOP

fornisce $\varphi' > \varphi$ sia $\widehat{\varphi} \vdash \varphi$
tale che $\varphi \rightarrow \neg\widehat{\varphi} \rightarrow \varphi'$

$\varphi(\alpha) \cup \{\widehat{\varphi}(\alpha)\} \vdash \perp$... ① ... $\neg\varphi(\alpha) \rightarrow \varphi'(\alpha)$ \square

Notazione $P^c - b_P(\varphi/A) = \{ \varphi(\alpha) \in \mathcal{F}^{P^c} : \not\vdash \varphi(\alpha) \}$

$\mathcal{H} - b_P(\varphi/A) = \{ \varphi(\alpha) \in \mathcal{H} : \not\vdash \varphi(\alpha) \}$

$b_P(\varphi/A) = \{ \varphi(\alpha) \in \mathcal{L} : \not\vdash \varphi(\alpha) \}$

Fatto $\varphi(\alpha) \subseteq \mathcal{F}^P(A)$ LSASE ① $\varphi(\alpha)$ è massimale compatibile in \mathcal{U}
② $\varphi(\alpha) = P^c - b_P(\varphi/A)$ per qualche $a \in \mathcal{U}$

Dim ① \Rightarrow ② prendiamo $a \models \varphi(\alpha)$

$$\varphi(\alpha) \subseteq P^c - b_P(\varphi)$$

\geq si ottiene dalle massimalità

③ Vedremo $\varphi(\alpha) \rightarrow b_P(\varphi/A)$
per $\varphi \in P$ i tipi di \mathcal{F}^P e \mathcal{L} sono gli stessi.

Dim ② \Rightarrow ① $\varphi(\alpha) \in \mathcal{F}^P(A) \nvdash \varphi$ allora $\neg\varphi(\alpha)$.

sempre da $\{\varphi(\alpha)\} \vdash \varphi(\alpha)$ allora $\neg\varphi'(\alpha)$ per qualche $\varphi' > \varphi$

ma $\varphi \vdash \varphi$ tale che $\varphi \rightarrow \neg\widehat{\varphi} \rightarrow \varphi'$ quindi $\widehat{\varphi}(\alpha)$

$\widehat{\varphi}(\alpha) \in P$ quindi $\varphi(\alpha) \rightarrow \neg\varphi(\alpha)$.

Definizione $\mathcal{M} \subseteq \mathcal{F}^{P^c}$ è P^c -denso se per ogni $\varphi' > \varphi \in \mathcal{F}^{P^c}$

c'è $\psi \in \mathcal{M}$ tale che $\varphi \rightarrow \psi \rightarrow \varphi'$

Proposizioni $\mathcal{M} \subseteq \mathcal{F}^{P^c}$ P^c -denso per ogni $\varphi \in \mathcal{F}^{P^c}$

(i) $\neg\varphi \rightarrow \bigvee \{ \neg\varphi_i \in \mathcal{M} : \neg\varphi_i \rightarrow \neg\varphi \text{ e } \varphi_i > \varphi \}$

(ii) $\neg\varphi \leftrightarrow \bigvee \{ \neg\varphi_i \in \mathcal{M} : \neg\varphi \rightarrow \neg\varphi_i \}$

Dimm $a \models \neg\varphi(\alpha)$ $\varphi(\alpha) = P^c - b_P(\varphi)$ è massimale

$$\begin{array}{c}
 p(\alpha) \longleftrightarrow p'(\alpha) \longleftrightarrow q(\alpha) = \text{fix-tp}(\alpha) \longleftrightarrow q'(\alpha) \\
 q'(\alpha) \longrightarrow \neg\varphi(\alpha) \qquad \qquad \neg\psi'(\alpha) \in \varphi' \text{ tells us } \neg\varphi(\alpha) \longrightarrow \neg\varphi(\alpha)
 \end{array}$$

$\varphi \longleftrightarrow \bigwedge \{ \psi \in \mathcal{A} : \varphi \rightarrow \psi \}$
 ↓ per def'n
 $\{ \psi \}'$

